SO That Looks Suspicious

Leveraging Process Memory & Kernel/Usermode Probes To Detect Shared Object Injection At Scale On Linux

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Whoami

Previously:

- ^o Senior security researcher @WithSecure/F-Secure.
- ^o Security research & endpoint agent developer @UKGov.
- ^o IR @Mandiant.
- ^o Prev Speaker @BlackHatUSA, BlackHatAsia x2 …

Professional interests:

- ^o OS internals.
- ^o Reverse engineering.
- ^o Tool & Sensor Dev.

OVERVIEW Shared Object Injection & the Linux threat landscape

O1 OVERVIEW Shared Object Injection & O2 ELF binary & process O3 HIJACKI ELF binary & process memory basics

Preloading and DT_NEEDED infections

AGENDA

01 **OVERVIEW**

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Shared Object Injection & The Linux Threat Landscape

(Trends)

The level of innovation of Linux malware came close to that of Windows-based malware, highlighting just how prevalent Linux malware innovation has become. a trend that we are sure to see increasing in 2022 as well.

IBM X-Force Threat Intelligence Index 2022

The importance of securing Linux® systems has risen in prominence as increasing amounts of malicious activity targeting Linux have appeared. Malware developers are increasingly developing Linux malware and creating Linux variants of existing malware families. These changes to the Linux threat landscape highlight the criticality of systems hardening and monitoring for malicious activity.

IBM X-Force Threat Intelligence Index 2024

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https://research.checkpoint.com/2023/the-platform-matters-a-comparative-studyon-linux-and-windows-ransomware-attacks/

https://vfeed.io/wp-content/uploads/2021/02/Top-10-Most-Used-MITRE-ATTCK.pdf

https://www.ncsc.nl/binaries/ncsc/documenten/publicaties/202 4/februari/6/mivd-aivd-advisory-coathanger-tlp-clear/TLP-CLEAR+MIVD+AIVD+Advisory+COATHANGER.pdf

■ Lack of detection maturity compared with Windows desktops.

Outdated Less detection and decedent \setminus Low visibility of $\quad \nearrow$ samples $\quad \diagup /$ research analysis

Threat Landscape

- Lack of detection maturity compared with Windows desktops.
- Threat groups incorporating opensource code directly into their malware: **at Landscape**

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alware:

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- Post exploitation frameworks & state sponsored attackers using SO injection techniques:

- Ninjasec/PupyRAT
- COATHANGER (Chinese FortiGate RAT)

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- Post exploitation frameworks & state sponsored attackers using SO injection techniques:
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- Open-source tooling & conferences presentations demonstrating Usermode memory injection techniques

 10 /

02 ELF 101

ELF Binary & Process Memory Basics

Binary Vs Memory Images

- ELF Sections contain large degree of forensic value.
	- Symbol Table, Relocation table, Constructors/Destructors, Program Data, Dynamic linking information

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Segments loose Section granularity!

Binary Vs Memory Images

- ELF Sections contain large degree of forensic value.
	- Symbol Table, Relocation table, Constructors/Destructors, Program Data, Dynamic linking information
- Segments loose Section granularity!
- Section header table is Optional in mapped memory image. Not suitable for use in forensic tooling.

Rebuilding Elf Sections From Memory

(Using The Dynamic Section)

Elf64_Word p_type; /* Segment type */ Elf64_Word p_flags; /* Segment flags */ Elf64 Off p offset; /* Segment file offset */ Elf64 Addr p vaddr; /* Segment virtual address */ Elf64_Addr p_paddr; /* Segment physical address */ Elf64 Xword p filesz; /* Segment size in file */ Elf64_Xword p_memsz; /* Segment size in memory */ Elf64 Xword p align; /* Segment alignment */

> 1. Start location of Dynamic Segment (PT_DYNAMIC) == 1:1 mapping of Dynamic

Rebuilding Elf Sections From Memory

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typedef struct

Elf64 Word p type: /* Segment type */ Elf64_Word p_flags; /* Segment flags */ Elf64 Off p offset; /* Segment file offset */ Elf64 Addr p vaddr; /* Segment virtual address */ Elf64_Addr p_paddr; /* Segment physical address */ Elf64 Xword p filesz: /* Seament size in file */ Elf64_Xword p_memsz; /* Segment size in memory */ Elf64 Xword p align: /* Segment alignment */ Elf64 Phdr;

> 1. Start location of Dynamic Segment (PT_DYNAMIC) == 1:1 mapping of Dynamic section.

2. Contains pointers to ELF sections needed by the Dynamic Linker

Readelf output of Dynamic section

root@test-VirtualBox:~# readelf -d /bin/bash

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03 DYNAMIC LINKER HIJACKING

Preloading & DT_NEEDED Infections

Abusing The Dynamic Linker To Load Malicious SOs

DT_NEEDED Entries & The Dynamic Linker

- Dynamic Segment 1:1 mapping with the Dynamic section
- Present in all dynamic linked binaries
- Each entry Dynamic section is required by the dynamic linker to load a binary into memory: **DT_NEEDED Entries & The Dy**

Mamic Segment 1:1 mapping with the Dynamic section

esent in all dynamic linked binaries

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esent in all dynamic linked binaries

contexts and a section

context in all dynamic linked binaries

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context in all dynam **DT_NEEDED ENTIRES & The Dynamic Segment 1:1 mapping with the Dynamic section**

esent in all dynamic linked binaries

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objec CONTRIGOT – Pointer should look to load libraries from.

CONTRIGOT – POINTED (SCIENCE TABLE (SOT).

ON THE DUMANNIC SECTION STEP (SOT).

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	-
	-
	-
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context-

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context-**DT_NEEDED ENTIRES & The Dynamic Segment 1:1 mapping with the Dynamic section**

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consecution

consecution and a binary into memory:

or DT_NEEDED - Dependencies to load.

or DT_NEEDED ○ DT_PLTGOT – Pointer within Global Offset Table (GOT).
	- DT_NEEDED Dependencies to load.
	-
	-
	-
	- the dynamic linker should look to load libraries from.

DT_NEEDED Infections (Overwrites)

Legitimate DT_NEEDED entry (libc.so)

Empty DT_DEBUG / DT_NULL entries

DT_NEEDED Infections (Overwrites)

apped into the ddress space

 \bullet

DT_NEEDED Infections

(Overwrites)

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7f69a9a92000-7f69a9a93000 r--p 0002c000 08:05 396631 7f69a9a93000-7f69a9a94000 rw-p 0002d000 08:05 396631 7f69a9a94000-7f69a9a95000 rw-p 00000000 00:00 0 7ffe91797000-7ffe917b8000 rw-p 00000000 00:00 0

ffffffffff600000-ffffffffff601000 --xp 00000000 00:00 0

7ffe917eb000-7ffe917ef000 r--p 00000000 00:00 0

7ffe917ef000-7ffe917f1000 r-xp 00000000 00:00 0

/home/vagrant/dt_infect/test_clean /home/vagrant/dt_infect/test_clean /home/vagrant/dt_infect/test_clean /home/vagrant/dt_infect/test_clean
/home/vagrant/dt_infect/test_clean [heap]

/usr/lib/x86_64-linux-gnu/libc-2.31.so $/usr/lib/x86_64-1inux-gnu/libc-2.31.so$ /usr/lib/x86_64-linux-gnu/libc-2.31.so /wsr/lib/x86_64-linux-gnu/libc-2.31.so
/wsr/lib/x86_64-linux-gnu/libc-2.31.so
/wsr/lib/x86_64-linux-gnu/libc-2.31.so
/wsr/lib/x86_64-linux-gnu/libc-2.31.so

/usr/lib/x86_64-linux-gnu/ld-2.31.so /usr/lib/x86 64 -linux-gnu/ld-2.31.so /usr/lib/x86_64-linux-gnu/ld-2.31.so
/usr/lib/x86_64-linux-gnu/ld-2.31.so
/usr/lib/x86_64-linux-gnu/ld-2.31.so

[stack] $[*var*]$ [vdso] [vsyscall]

 $, 61$

DT_NEEDED Infections

(Overwrites)

DT_NEEDED Infections (Insertions)

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DT_NEEDED Infections (Insertions)

 \bullet

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DT_NEEDED Infections (Detection)

x OVERWRITES

- o Order of DT_NEEDED entries in dynamic section
- o Dynamic string table extension
- o Missing DT_DEBUG/DT_NULL entries
- o Header manipulation

INSERTIONS

- o Evidence of Program header relocation
- o Dynamic string table extension
- o Does SO name ptr point within dynamic string table.
- o Duplication of Symbol names across Shared Objects

 \circ

DT_NEEDED Infections (Detection)

x OVERWRITES

},

"dt needed index": 0, "index_into_strtab": 1, "module_name": "libc.so.6", "name_in_dynstr": true

"dt_needed_index": 12, "index into strtab": 68, "module name": "libevil.so", "name in dynstr": false

"dt_needed_wrong_order": true, "dt_null_present": true, "debug section present": false, "dynstr_manipulated": true, "headers manipulated": true,

 \cap

Preloading Abuse (LD_PRELOAD)

- Preloaded SO functions overwrite functions of non-preloaded SOs. Acting as a search order hijacking mechanism.
- Preloading mechanisms:
	- o LD_PRELOAD env var
	- Dynamic linker '—preload' flag
	- /etc/ld.so.preload
- Preloading has legitimate uses: for debugging / compatibility
- Offers attackers a simple way to install hooks / execute constructor code
- Used by:
	- Azazel, BEURK, Jynx, Vlany. Umbreon Usermode rootkits
	- HiddenWasp malware & Threat groups.

The Problem With Detecting Malicious Preloading

- Current detection solutions only monitor 'existence' of preloading rather than 'effect':
	- Command lines, paths & env variables.
	- Still requires manual analysis

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- Identify the individual hooks?
- Which preloaded SOs are responsible?
- Where is the location of the hook in memory?
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The Problem With Detecting Malicious Preloading

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Step 1: Identifying All Imports **Step 1: Identifying All Imports**
i. ELF executable header fields e*_phoff & e_phentsize ->* The
program header table.

program header table.

- program header table.
- segment.

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- iii. Dynamic segment is 1:1 mapping of the *dynamic section* $\begin{bmatrix} 0 \\ 1 \end{bmatrix}$ containing pointers to:
	- Global offset table (DT_PLTGOT).
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	- Dynamic symbol table (DT_SYMTAB).
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	- The address of the associated **GOT** entry
	- relates to (Elf64 Sym).

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	- The address of the associated **GOT** entry
	- relates to (Elf64 Sym).
- VI. Elf64 Sym entry contains the offset within the dynamic string

Step 2: Establish A List Of SOs & Their Base Addresses

- i. Locate following sections in .dynamic section:
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	- Debug section (DT_DEBUG).

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	- \blacksquare Using the GOT got[1]
- iii. Iterate through the link map linked list and extract the loaded base address for each SO in memory.

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Identify preloaded SOs:

- Reading LD_PRELAOD from the stack.
- Reading /etc/ld.so.preload.

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	- Dynamic symbol table (DT_SYMTAB).
	- **Dynamic string table (DT STRTAB).**
- ii. Determine the number of symbol table entries using either:
	- Hash table.
	- GNU Hash table.
- iii. Only collect exported symbols which are:
	-
	-

Step 4 & 5: Comparisons & Matching Symbol Names

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■ Compare imported symbols with exported symbols from the any preloaded SOs.

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- Match legitimate export names with names of hooks to identify victim SOs.

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Step 5:

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Malicious hook code @VA

ELFieScanner

- Linux process memory scanning tool that detects various forms of:
	- Shared Object injection.
	- Shellcode injection & Process hollowing.
	- Entry point manipulation.
	- API Hooking.
- 43 different heuristics, controllable via configuration file.
- Multithreaded, written in C++, scans both x86/x64 processes.
- Outputs data into NDJSON file
- https://github.com/JanielDary/ELFieScanner

04 REALTIME INJECTION & TARGETING

57 \degree Reflective Shared Object Injection & __libc_dlopen_mode()

Attack Techniques

Existing Real-time Detection Strategies (SO Injection)

Solutions include:

- Monitoring/restricting the use of PTRACE() syscalls.
- Enumerating /proc/<pid>/maps file for RWX regions.
- Combining output with file events and command lines on a best effort basis.
- Blindly scanning memory with Yara signatures.

Issues with current solutions:

○ Browsers, debuggers, AVs and interpreters can exhibit the same behaviors in a legitimate way.

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- Solutions do not target SO injection specifically.
- Can introduce data volume and backend performance issues.
- Lots of data for an analyst to sift through.

- Two functions can be used to load a SO into a Linux process:
	- dlopen().
	- __libc_dlopen_mode().
- Force a victim process to call either function ensures the dynamic linker does most of the work.
	- __libc_dlopen_mode() almost always targeted over dlopen().

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- Method 1: Writing own injector:
	- Attach to a victim process with LIBC loaded.
	- Resolve the address of __libc_dlopen_mode() & modify the instruction pointer.
	- Replace registers (x64) or stack values (x86) with the correct arguments.
	- Resume execution.

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	- Resume execution.
- Method 2: Using a GDB bash one-liner.

echo 'print __libc_dlopen_mode("/tmp/sample_library.so", 2)' | gdb -p <PID>

Monitoring Function Calls

RA

- **LTrace()** \bullet **l** \bullet Can only target individual / groups of processes.
	-
	- o Malicious processes can prevent itself being debugged using PTRACE_TRACEME

Monitoring Function Calls

o Can only target individual / groups of processes.

o Uses PTRACE – (Slow + Invasive)

o Malicious processes can prevent itself being debugged using PTRACE_TRAC Monitoring Function Calls

|小||

- **LTrace()** $\bigcap_{i=1}^{\infty}$ **Can only target individual / groups of processes.**
	-
	- o Malicious processes can prevent itself being debugged using PTRACE_TRACEME

Dynamic Instrumentation (DI)

Uprobes

- o Introduced in Linux 3.5
- o System wide effect
- o DI of User level functions & offsets
- o Defining a Uprobe requires:
	- Path of SO.
	- Offset to target function
	- Selected function parameters & corresponding register/stack values.

Kprobes

- o Introduced in Linux 2.69
- o System wide effect
- o DI of Kernel level functions & offsets

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o Often used by eBPF programs

Calculating The Offset To __libc_dlopen_mode (Uprobe)

■ Method 1: Using nm -

nm -D /usr/lib/x86_64-linux-gnu/libc-2.32.so | grep __libc_dlopen_mode

Calculating The Offset To __libc_dlopen_mode (Uprobe) **Calculating The Offset To**
(Uprobe)
Method 1: Using nm -

nm -D /usr/lib/x86_64-linux-gnu/libc-2.32.so | grep _libc_dlopen_mod
Method 2: Manually enumerating the SO on disk:

i. Locate Section Hdrs table via Exe Hdrs.

- Method 2: Manually enumerating the SO on disk:
	-

Calculating The Offset To __libc_dlopen_mode (Uprobe) **Calculating The Offset To**
(Uprobe)
Method 1: Using nm -

nm -D /usr/1ib/x86_64-1inux-gnu/1ibc-2.32.so | grep _1ibc_dlopen_mod
Method 2: Manually enumerating the SO on disk:

i. Locate Section Hdrs table via Exe Hdrs.

ii **Calculating The Offset To**

(Uprobe)

Method 1: Using nm -

nm -D /usr/1ib/x86_64-1inux-gnu/1ibc-2.32.so | grep _1ibc_dlopen_mod

Method 2: Manually enumerating the SO on disk:

i. Locate Section Hdrs table to locate:

ii

- Method 2: Manually enumerating the SO on disk:
	-
	- - the dynamic symbol table (.dynsym)
		- dynamic string table (.dynstr)

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- Method 2: Manually enumerating the SO on disk:
	-
	- - the dynamic symbol table (.dynsym)
		- dynamic string table (.dynstr)
- **iii.** Enumerate .dynsym 8. dynstr tables to match symbol and the symbol and the symbol dable (.dynsym)

iii. Enumer names with Elfxx_Sym entry.

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ii Calculating The Offset To __libc_dlopen_mode (Uprobe)

- Method 2: Manually enumerating the SO on disk:
	-
	- - the dynamic symbol table (.dynsym)
		- dynamic string table (.dynstr)
- **iii.** Il is Simple . The dynsym and the symbol of the symbols of the dynamic symbol table (dynsym)

ii. Use Section Hdrs table to locate:
 iii. Use Section Hd names with Elfxx_Sym entry. Method 1: Using nm -

nm -D /usr/1ib/x86_64-1inux-gnu/1ibc-2.32.so | grep __libc_dlopen_mode

Method 2: Manually enumerating the SO on disk:

i. Locate Section Hdrs table via Exe Hdrs.

ii. Use Section Hdrs table via Exe H
	- determine its file offset.

Determining The Function Parameters (Uprobe) **Example 12 Constraining The Function

Frameters (Uprobe)**

libc_dlopen_mode() uses same two parameters as dlopen():
 Allocaters - Path of SO (rdi).

Final prode - Loading method flag (rsi).

Lentify any variations betwe

- _libc_dlopen_mode() uses same two parameters as dlopen():
	-
	-
- Identify any variations between GLIBC versions.

```
C di-libele 2 \timesFunction prototype
elf > C di-libc.c > @ _libc_diopen_mode(const char *
     f^* ... and these functions call dle
      \sqrt{010} *
       _libc_dlopen_mode (const char *name, int mode)
        struct do dlopen args args;
        args.name = name;
        args.mode = mode;args.caller_dlopen = RETURN_ADDRESS (0);
156
      #ifdef SHARED
        if (irtid_active ())
 160
         return GLRO (dl dlfcn hook)->libc dlopen mode (name, mode);
      #endif
 162
        return dlerror_run (do_dlopen, &args) ? NULL : (void *) args.map;
 163
      \left| \right\rangleC difcn.h \timesusr > include > x86 64-linux-gnu > bits > C difcn.h > ...
     /* The MODE argument to 'dlopen' contains one of the following: */
     #define RTLD_LAZY 0x00001 /* Lazy function call binding. */
      #define RTLD_NOW 0x00002 /* Immediate function call binding. */
      #define RTLD BINDING MASK 0x3 /* Mask of binding time value. */
 -26#define RTLD_NOLOAD 0x00004 /* Do not load the object. */
      /* If the following bit is set in the MODE argument to 'dlopen',
         the symbols of the loaded object and its dependencies are made
         visible as if the object were linked directly into the program. */
     #define RTLD GLOBAL 0x00100
      /* Unix98 demands the following flag which is the inverse to RTLD_GLOBAL.
         The implementation does this by defau
 36
         value to zero. */
                                                    Loading method 
      #define RTLD_LOCAL 0
      /* Do not delete object when closed.
 48
                                                     flags#define RTLD NODELETE 0x01000
```


Determining The Function Parameters (Uprobe) **Example 12 Constraining The Function

Frameters (Uprobe)**

libc_dlopen_mode() uses same two parameters as dlopen():
 Allocaters - Path of SO (rdi).

Final prode - Loading method flag (rsi).

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	- The first parameter 'path' renaming this to 'injected lib' from the rdi register.

Brendan Gregg's F-Trace Uprobe wrapper:

./uprobe -H 'p:/lib/x86_64-linux-gnu/libc-2.32.so:0x1598a0 injected_lib=+0(%di):string

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format.

Brendan Gregg's F-Trace Uprobe wrapper:

./uprobe -H 'p:/lib/x86_64-linux-gnu/libc-2.32.so:0x1598a0 injected_lib=+0(%di):string mode=%si:x32'

Methods Of Detecting The Injector Process

1. Using existing telemetry to find the most recent PTRACE_ATTACH event prior to the Uprobe firing. This will be the injector process

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- 2. Signature on command line arguments supplied to GDB containing '*_libc_dlopen_mode'*.
- 3. Search a running process' .rodata section for references to __libc_dlopen_mode():
	- Only works if the injector process still exists.

- The Linux equivilent of Reflective DLL injection on Windows, used by:
	- **InfoSecguerrilla/ReflectiveSOInjection tool.** infosecguerrilla / ReflectiveSOInjection Ω
	- **N1nj4sec/Pupy framework.** n1nj4sec / pupy
- Facilitates the loading of a SO directly from memory by using a custom loader:
	- **Allocates a RWX anonymous memory region.**
	- **Maps a SO into the region.**
	- Uses Libc exports to resolve symbols and perform relocations.

● Current detection strategies rely on identifying existing RWX regions, this can be easily circumvented by:

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Using Kprobes To Target Memory Allocations

- Target the initial memory allocation.
- Exported Kernel Symbols found in /proc/kallsyms.
- mmap() not exported:
	- Internally calls sys_mmap->ksys_mmap_pgoff.

```
148
       asmlinkage unsigned long
149
       sys mmap (unsigned long addr, unsigned long len, int prot, int flags, int fd, long off)
150
       \{151
               if (offset_in-page(off) != 0)152
                        return -EINVAL;
153
154
               addr = ksys_mmap_pgoff(addr, len, prot, flags, fd, off >> PAGE_SHIFT);
155
               if (!IS ERR((void * ) addr))156
                        force successful syscall return();
157
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```


ifdef MAP ANONYMOUS

define MAP_ANONYMOUS 0x20

else

endif

/* Don't use a file. $*/$

The Probe & Telemetry

./kprobe 'p:ksys mmap pgoff addr=%di:x32 len=%si:x32 prot=%dx:x32 flags=%cx:x32 fd=%r8:x32 off=%r9:x32' 'flags== $0x2288prot=-0x7'$

root@ubMalware:~/perf-tools/kernel# ./kprobe 'p:ksys_mmap_pgoff addr=%di:x32_len=%si:x32_prot=%dx:x32_flags=%cx:x32_fd=%r8:x32_off=%r9:x32' 'flags==0x22&&prot==0x7 Tracing kprobe ksys mmap pgoff. Ctrl-C to end. victim process-6030

- A Kprobe can be used to target:
	- Anonymous memory allocations.
	- With initial RWX / RX permissions.
- Multiple probes can be set for each allocation variation & change e.g. mprotect()

The Probe & Telemetry

./kprobe 'p:ksys mmap pgoff addr=%di:x32 len=%si:x32 prot=%dx:x32 flags=%cx:x32 fd=%r8:x32 off=%r9:x32' 'flags==0x22&&prot==0x7'

root@ubMalware:~/perf-tools/kernel# ./kprobe 'p:ksys_mmap_pgoff addr=%di:x32 len=%si:x32 prot=%dx:x32 flags=%cx:x32 fd=%r8:x32 off=%r9:x32 'flags==0x22&&prot==0x7 Tracing kprobe ksys mmap pgoff. Ctrl-C to end. victim process-6030 [001] 10682.995045: ksys_mmap_pgoff: (ksys_mmap_pgoff+0x0/0x2a0) addr=0x0 len=0xc930 prot=0x7 flags=0x22 fd=0xffffffff off=0x0

- A Kprobe can be used to target:
	- Anonymous memory allocations.
	- With initial RWX / RX permissions.
- Multiple probes can be set for each allocation variation & change e.g. mprotect()
- Capture the memory address & length supplied to ksys_mmap_pgoff to trigger a targeted scan.

05 HIDE & SEEK

Hidden Shared Objects & Detection Rules

Hidden Shared Objects

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Process Mappings

The 'proc/<pid>/maps' is the pseudo-filesystem representation of a process' memory mappings, this includes it's loaded SOs

Monero miner (libprocesshider)

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Hidden Shared Objects

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Monero miner (libprocesshider)

Hidden Shared Objects

(Enumeration Methods)

Process Mappings

The 'proc/<pid>/maps' is the pseudo-filesystem representation of a process' memory mappings, this includes it's loaded SOs

GOT[1] / DT_DEBUG Contains the address of the link_map structure linked list, containing the base address & name of loaded SO's

link_map

DT_NEEDED

.Dynamic Section DT_NEEDED entry type contains names of SOs to load at runtime via standard search order mechanisms.

Hidden Shared Objects (Rules) **2.** SOs that only appear in either the *link_map* OR *proc/*<*pid>/maps* <u>but not both!</u>
2. SOs with the <u>same name but different base addresses</u> in *proc/<pid>/maps*. 4. DT_NEEDED entries <u>missing</u> from either the *link_*

- SOs that only appear in either the link_map OR proc/<pid>/maps but not both!
-
-
-
- 5. SOs with non-standard paths.

Cheat Sheet

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KEY TAKEAWAYS

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Key Takeaways

1.

Less spotlight on the Linux threat landscape leading to lower detection maturity when compared to Windows

Telemetry & tooling needs to be kept up to date otherwise simple modifications can sidestep existing rules.

2.

3.

Utilizing K/Uprobes as targeted triggers can greatly reduce performance overheads when running memory scanners, opening up their applicable use cases.

